

Function approximation with one-bit Bernstein and neural networks

Weilin Li

Courant Institute of Mathematical Sciences

weilinli@cims.nyu.edu

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Figure: Sinan Güntürk

References:

- Approximation with one-bit polynomials in Bernstein form.
arXiv:2112.09183
- Approximation of functions with one-bit neural networks.
arXiv:2112.09181

Can also be found on my webpage:

weilinli@cims.nyu.edu

1 Introduction

2 Why Bernstein?

3 Back to Neural Networks

4 Main Results

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Main question

Can we approximate any reasonable function with a **coarsely quantized** neural network?

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- A feed forward neural network is a function represented as a composition,

$$x \mapsto \rho_L(A_L(\cdots \rho_1(A_1x + b_1)) + b_L),$$

where each weight matrix A_ℓ and bias vector b_ℓ are real.

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- A **quantized** NN is one where all non-zero parameters (weights and biases) must be selected from a set of discrete values (“alphabet”).

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- A **coarsely quantized** NN uses an alphabet with a small collection of values. For example, a one-bit alphabet,

$$\mathcal{A}_1 := \{\pm 1\}.$$

Motivation 1: Practical issue

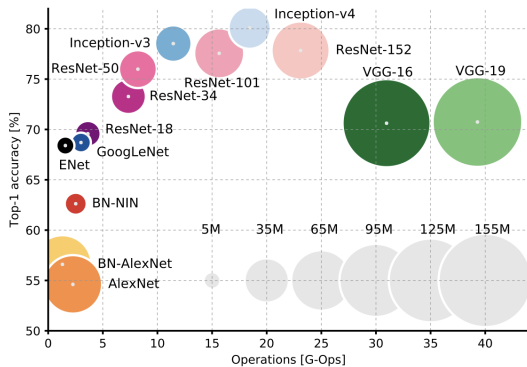


Figure: Memory footprint of some networks [Canziani, Paszke, Culurciello 16]

How to transfer huge networks onto low power devices (phones, cars, tablets, etc.)?

Survey article [Yunhui Guo 18]

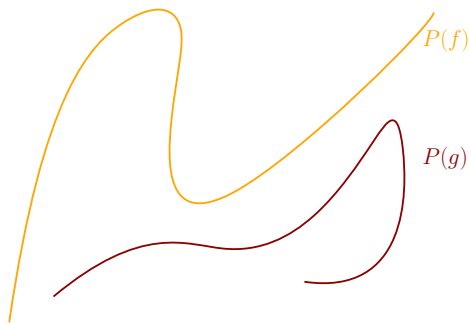
- Quantize the network → low-bit representation
- cheaper storage cost and faster forward pass
- less energy consumption

Use of sophisticated quantization techniques: [Lybrand, Saab 21], [Ashbrock, Powell 21]

Motivation 2: Over-parameterization

For each f , look at the parameters that represent f ,

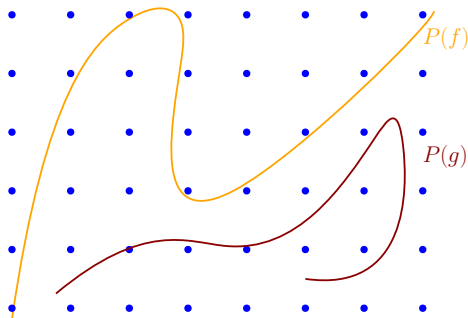
$$P(f) = \{\theta \in \mathbb{R}^D : NN_{\theta} = f\}$$



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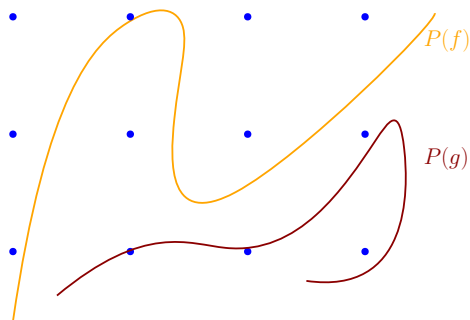
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What makes the main question difficult to answer?

Numerous results for approximation by neural networks (incomplete list of course):

[Cybenko 89], [Barron 94], [Yarotsky 17], [Shaham, Cloninger, Coifman 18], [Bolcskei, Grohs, Kutyniok, Petersen 19], [Daubechies, DeVore, Foucart, Hanin 21], [Lu, Shen, Yang, Zhang 21]

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Main strategy of the above approximation strategies:

For fixed f ,

$$f \approx \sum_k a_k \phi_k$$

- $\{\phi_k\}_k$ are implementable by a neural network with carefully chosen parameters, e.g., local polynomials, ridge functions, wavelets, Fourier
- $\{a_k\}_k \subseteq \mathbb{R}$ depend on f , e.g., basis or frame coefficients

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Approach needs to be modified when we consider coarsely quantized networks!

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Find a suitable set of functions $\{\phi_k\}_k$ such that:

- **Approximation.** Linear combinations of $\{\phi_k\}_k$ can efficiently approximate large function classes.
- **Implementation.** The $\{\phi_k\}_k$ can be implemented by a **coarsely quantized** neural network.

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- **Quantization.** Coefficients in the $\{\phi_k\}_k$ basis can be quantized in an effective way. For each $\{a_k\}_k$ we can find $\{\sigma_k\} \subseteq \mathcal{A}$ such that

$$\sum_k a_k \phi_k - \sum_k \sigma_k \phi_k \text{ is small.}$$

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Three term decomposition:

$$f - f_{NN} = \underbrace{\left(f - \sum_k a_k \phi_k \right)}_{\text{approximation by } \phi \text{ error}} + \underbrace{\left(\sum_k a_k \phi_k - \sum_k \sigma_k \phi_k \right)}_{\text{quantization error}} + \underbrace{\left(\sum_k \sigma_k \phi_k - f_{NN} \right)}_{\text{implementation error}}.$$

Kantorovich: If $f: [0, 1] \rightarrow \mathbb{R}$ is continuous and $f(0)$ and $f(1)$ are integers, then the function

$$B_n^*(f)(x) := \sum_{k=0}^n \left[f\left(\frac{k}{n}\right) \binom{n}{k} \right] x^k (1-x)^{n-k},$$

converges uniformly to f as $n \rightarrow \infty$. Here, $[t]$ is rounding t to the nearest integer.

Proof can be found in Chapter 2.4 of *Constructive Approximation: Advanced Problems* by Lorentz, v Golitschek, and Makovoz.

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- Implies that polynomials with integer coefficients are dense in the space of continuous functions.
- Main drawback is that if $\|f\|_\infty \leq 1$, then

$$\left| \left[f\left(\frac{k}{n}\right) \binom{n}{k} \right] \right| \text{ could be as big as } 2^n,$$

so not helpful for coarse quantization.

Bernstein polynomials of degree n : for each $0 \leq k \leq n$,

$$p_{n,k}(x) := \binom{n}{k} x^k (1-x)^{n-k}.$$

Berstein: For any $f \in C([0, 1])$,

$$\|f - B_n(f)\| = \left\| f - \sum_{k=0}^n f\left(\frac{k}{n}\right) p_{n,k} \right\|_{\infty} \rightarrow 0.$$

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Question: For any $f \in C([0, 1])$ and $\epsilon > 0$, is there an integer n and $\{\sigma_k\}_{k=0}^n$ such that $\sigma_k \in \{\pm 1\}$ and

$$\left\| f - \sum_{k=0}^n \sigma_k p_{n,k} \right\|_{\infty} \leq \epsilon?$$

Bernstein polynomials of degree n : for each $0 \leq k \leq n$,

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Question: For any $f \in C([0, 1])$ such that $\|f\|_\infty \leq 1$ and $\epsilon > 0$, is there an integer n and $\{\sigma_k\}_{k=0}^n$ such that $\sigma_k \in \{\pm 1\}$ and for all x away from the endpoints,

$$\left| f(x) - \sum_{k=0}^n \sigma_k p_{n,k}(x) \right| < \epsilon?$$

Theorem

For any Lipschitz continuous f with $\|f\|_\infty \leq 1$ and integer n , there exist $\{\sigma_k\}_{k=0}^n \subset \{\pm 1\}$, such that

$$\left| f(x) - \sum_{k=0}^n \sigma_k p_{n,k}(x) \right| \lesssim \frac{1 + |f|_{Lip}}{1 + \sqrt{nx(1-x)}}.$$

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- By density, can extend to continuous functions.
- In comparison to Kantorovich, our construction uses significantly fewer bits.
- Implies a L^p estimate: for all $p \in [1, 2)$,

$$\left\| f - \sum_{k=0}^n \sigma_k p_{n,k} \right\|_{L^p([0,1])} \lesssim \frac{1 + |f|_{Lip}}{\sqrt{n}}.$$

A constructive proof

The sequence of ± 1 are chosen through an algorithm called first order $\Sigma\Delta$. Set $a_k = f(k/n)$ and we find σ such that it satisfies a finite difference equation

$$a_k - \sigma_k = (\Delta u)_k := u_k - u_{k-1}.$$

Can always be done with a $\|u\|_\infty \lesssim 1$ due to the assumption that $\|f\|_\infty \leq 1$. Explicitly,

$$u_k := u_{k-1} + a_k - \sigma_k, \quad \sigma_k := \text{sign}(a_k + u_{k-1}).$$

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Bound the total error:

$$\begin{aligned} & \left| f(x) - \sum_{k=0}^n \sigma_k p_{n,k}(x) \right| \\ & \leq \left| f(x) - \sum_{k=0}^n a_k p_{n,k}(x) \right| + \left| \sum_{k=0}^n (a_k - \sigma_k) p_{n,k}(x) \right| \\ & \lesssim \frac{|f|_{Lip}}{\sqrt{n}} + \left| \sum_{k=0}^n (\Delta u)_k p_{n,k}(x) \right| \end{aligned}$$

Defining $p_{n,k+2} := 0$,

$$\begin{aligned} \left| \sum_{k=0}^n (\Delta u)_k p_{n,k}(x) \right| &= \left| \sum_{k=0}^n u_k (p_{n,k}(x) - p_{n,k+1}(x)) \right| \\ &\leq \|u\|_\infty \sum_{k=0}^n |p_{n,k}(x) - p_{n,k+1}(x)| \quad (\text{Frame variation}) \\ &= \frac{\|u\|_\infty}{(n+1)x(1-x)} \sum_{k=0}^n |(k+1) - (n+1)x| p_{n+1,k+1}(x) \\ &= \frac{\|u\|_\infty}{(n+1)x(1-x)} \left(\sum_{k=0}^n |(k+1) - (n+1)x|^2 p_{n+1,k+1}(x) \right)^{1/2} \\ &\lesssim \frac{1}{\sqrt{nx(1-x)}}. \end{aligned}$$

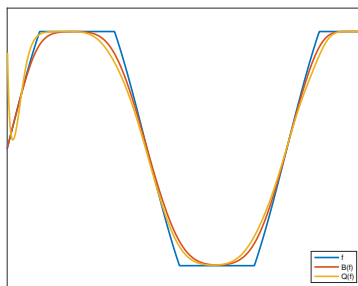
Frame variation 2 papers, [Benedetto, Powell, Yilmaz, 06]

Simple demonstration

Consider the continuous function,

$$f(x) = \max\left(\min\left(\frac{3}{2}\sin(8x), 1\right), -1\right).$$

Plot of $B_n(f)$ and constructed quantized approximation, for $n = 100$.



Coefficients of quantized approximation (where $-1 \leftrightarrow 0$):

```
10101111111111111111111111111111110111101010010000000  
00000000000000000000000010001001101110111111111111111
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The previous results give us hope that Bernstein polynomials satisfy our “wish list”, but there is still a main concern when applying these results to networks:

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high dimensionality \longrightarrow spaces of smooth functions.

Multivariate Bernstein of order n : For each $k \in \mathbb{N}^d$ with $0 \leq k_j \leq n$,

$$p_{n,k}(x) := p_{n,k_1}(x_1) \cdots p_{n,k_d}(x_d).$$

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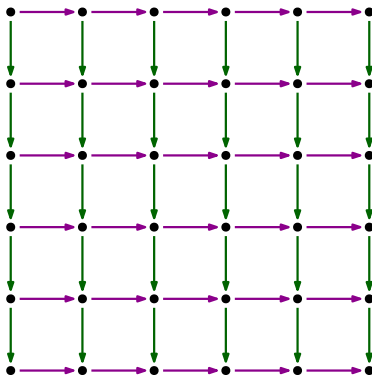
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- **High dimensionality** is a major issue for $\Sigma\Delta$. No available $O(1)$ -alphabet result for stable d dimensional $\Sigma\Delta$.
- **Higher order smoothness** is an issue do to this saturation result:
If $f \in C^2([0, 1])$, then

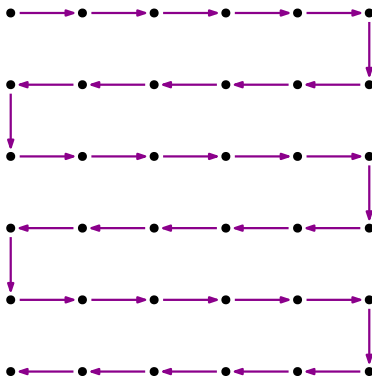
$$\lim_{n \rightarrow \infty} n(B_n(f)(x) - f(x)) = \frac{f''(x)x(1-x)}{2}.$$

Bernstein polynomial of f is unable to exploit additional smoothness.

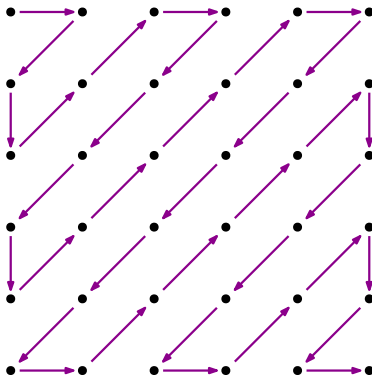
Say $d = 2$. How to order $\{p_{n,k}\}_{k_1=0,k_2=0}^n$?



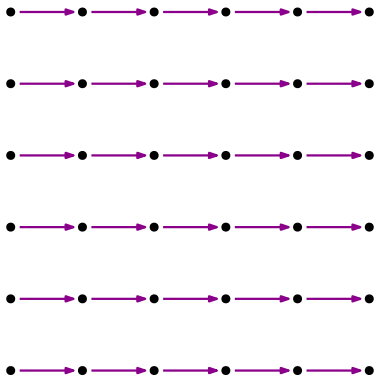
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Theorem

Pick any $1 \leq \ell \leq d$ and for any $\{a_k\}_{0 \leq k \leq n}$ with $\|a\|_\infty \leq \mu < 1$, let $\{\sigma_k\}_{0 \leq k \leq d}$ be a solution to a r -th order $\Sigma\Delta$ scheme applied only in the ℓ -th direction,

$$a - \sigma = (\Delta_\ell)^r u.$$

Then we have

$$\left| \sum_{0 \leq k \leq n} a_k p_{n,k}(x) - \sum_{0 \leq k \leq n} \sigma_k p_{n,k}(x) \right| \lesssim_{r,\mu} \min \left(1, n^{-r/2} x_\ell^{-r} (1 - x_\ell)^{-r} \right).$$

Relies on stability of one-dimension r -th order $\Sigma\Delta$ schemes

[Daubechies, Devore 03], [Güntürk 03].

Saturation: If $f \in C^2([0, 1])$, then

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How to get around the saturation problem?

Saturation: If $f \in C^2([0, 1])$, then

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How to get around the saturation problem?

Classical work: [Micchelli 73], [Felbecker 79] Define the iterated Bernstein operator

$$U_{n,m}(f) := \left(I - (I - B_n)^m \right)(f).$$

The for every $f \in C^s([0, 1])$,

$$\|f - U_{n, \lceil s/2 \rceil}(f)\|_\infty \lesssim_s \|f\|_{C^s} n^{-s/2}.$$

Takes advantage of smoothness, but is $U_{m,n}(f)$ a linear combination of Bernstein polynomials with small enough coefficients (necessary for quantization)?

Theorem

Fix any integers $d, s \geq 1$, $\delta \in (0, 1)$, and $f \in C^s([0, 1]^d)$. Then for all integers

$$n \geq \frac{sd^2 \|f\|_{C^2}}{4\delta},$$

there exist $\{a_k\}_{0 \leq k \leq n}$ such that

$$\left\| f - \sum_{0 \leq k \leq n} a_k p_{n,k} \right\|_{\infty} \lesssim_{s,d} \|f\|_{C^s} n^{-s/2},$$

and

$$\|a\|_{\infty} \leq \|f\|_{\infty} + \delta.$$

How to implement Bernstein polynomials?

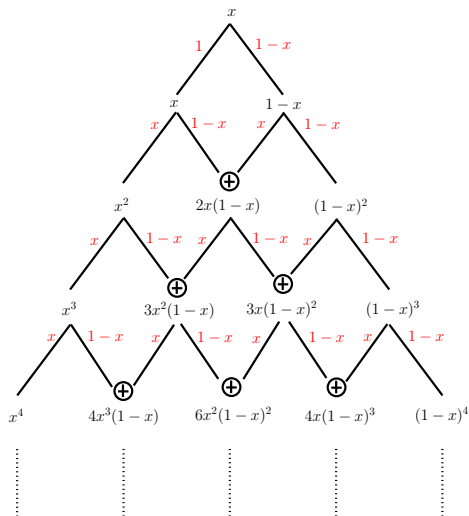


Figure: Pascal triangle implementation of Bernstein

Key formula:

$$ab = \frac{1}{2}(a + b)^2 - \frac{1}{2}a^2 - \frac{1}{2}b^2$$

Need to find a coarsely quantized neural network that can implement squaring.
Suffices for only non-negative a, b .

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- Easy if we use the quadratic non-linearity, $\rho(t) = t^2/2$, so

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- Squaring can be efficiently approximated by a ReLU network [Yarotsky 17], by exploiting the formula

$$x(1-x) = \sum_{m=1}^{\infty} \frac{\phi^{om}}{4^m}, \quad \phi(x) = \begin{cases} 2x & \text{if } 0 \leq x \leq \frac{1}{2}, \\ 2-2x & \text{if } \frac{1}{2} \leq x \leq 1. \end{cases}$$

Can do approximate multiplication with a $\{\pm 1/2, \pm 1, \pm 2\}$ -quantized ReLU network.

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$$f - f_{NN} = \underbrace{f - f_B}_{\text{Bern. approx. error}} + \underbrace{f_B - f_Q}_{\text{Bern. quan. error}} + \underbrace{f_Q - f_{NN}}_{\text{Bern. implementation error}} .$$

- f_B = iterated Bernstein of f , then covered to linear Bernstein.
- f_Q = function obtained from r -th order $\Sigma\Delta$ applied to coefficients of f_B .
- f_{NN} = coarsely-quantized neural network approximation of f_Q .

Theorem

Fix any integers $s, d \geq 1$, $\mu \in (0, 1)$, and any $f \in C^s([0, 1]^d)$ with $\|f\|_\infty \leq \mu$.

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$$n \geq \frac{sd^2 \|f\|_{C^2}}{2(1 - \mu)},$$

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$$|f(x) - f_{NN}(x)| \lesssim_{s,d,\mu} \|f\|_{C^s} \min(1, n^{-s/2} x_\ell^{-s} (1 - x_\ell)^{-s}).$$

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f_{NN} can be chosen as either

- a $\{\pm 1\}$ -quantized quadratic neural network with

$$O(n + d) \text{ layers,}$$

$$O(dn^2 + dn^d) \text{ nodes and parameters.}$$

- a $\{\pm 1/2, \pm 1, \pm 2\}$ -quantized ReLU neural network with

$$O\left((n + d)(d + s/2) \log n\right) \text{ layers,}$$

$$O\left((n^2 + dn^d)(d + s/2) \log n\right) \text{ nodes and parameters.}$$

Network in the main theorem

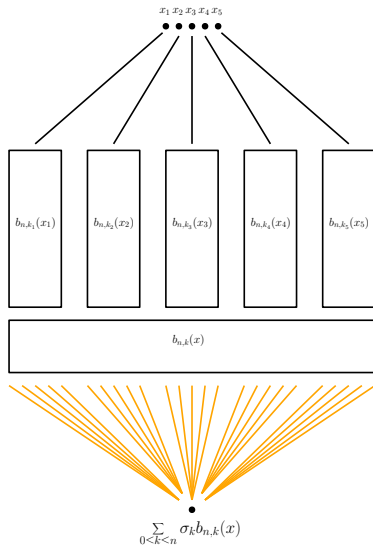


Figure: Coarsely-quantized network

Binary Bernstein Algorithm

System parameters and assumptions:

- d , dimension or number of variables,
- s , smoothness of the target function,
- $\mu \in (0, 1)$, upper bound on $\|f\|_\infty$,
- n , any integer at least $sd^2\|f\|_\infty^2/(2(1-\mu))$.

Algorithm:

- Input samples $\{f(k/n)\}_{0 \leq k \leq n}$.
- Calculate $\{a_k\}_{0 \leq k \leq n}$ defined to be

$$a_k := U_{n, \lceil s/2 \rceil}(f)\left(\frac{k}{n}\right).$$

- Run s -th order $\Sigma\Delta$ quantization in the ℓ -th direction on $\{a_k\}_{0 \leq k \leq n}$ to obtain

$$\{\sigma_k\}_{0 \leq k \leq n} \subset \{\pm 1\}.$$

- Guaranteed that

$$\left| f(x) - \sum_{0 \leq k \leq n} \sigma_k p_{n,k}(x) \right| \lesssim_{s,d,\mu} \|f\|_{C^s} \min(1, n^{-s/2} x_\ell^{-s} (1-x_\ell)^{-s}).$$

Can these results be adapted and generalized to other approximation systems?

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Some properties that we used:

- Bernstein polynomials are non-negative and form a partition of unity.
- Growth conditions on central moments.
- Bernstein polynomials are localized but not too much.
- Explicit formulas for $p_{n,k+1} - p_{n,k}$; in particular, there is significant cancellation.

Given a trained neural network F we can query it at lattice points $\{F(k/n)\}_{0 \leq k \leq n}$ and use our algorithm to produce a one-bit coefficients

$$\{\sigma_k\}_{0 \leq k \leq n} \subseteq \{\pm 1\}.$$

(Can even be computed in parallel.)



Figure: Sinan Güntürk

References:

- Approximation with one-bit polynomials in Bernstein form.
arXiv:2112.09183
- Approximation of functions with one-bit neural networks.
arXiv:2112.09181

Can also be found on my webpage:

weilinli@cims.nyu.edu

Thank you!!